

CSc 461/561

Multimedia Systems

Multimedia Congestion Control

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What we have so far

- Multicast
 - IP multicast, IGMP, multicast routing
- Resource reservation
 - RSVP, IntServ, DiffServ
- Transport
 - RTP and RTCP
- Signaling
 - SIP

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What still missing

- End-to-end flow, error, congestion control
 - how to coordinate sender, receiver, network
 - how to detect, correct, recover errors
- For TCP/IP-based data applications
 - TCP flow, error, congestion control
 - embedded and integrated
- For multimedia applications
 - TCP is not the best choice!

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Multimedia traffic control

- Flow, error, congestion control
 - should be separate, flexible, adaptive
 - due to diverse application requirements
- Flow control
 - window control, rate control, bucket-controlled
- Error control
 - redundancy, forward correction, retransmission
- Congestion control

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TCP-friendly congestion control

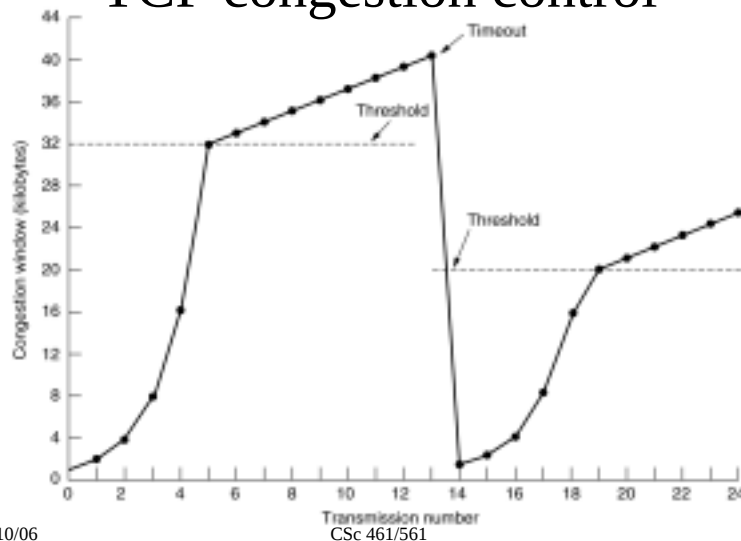
- What is TCP-friendly?
 - in a comparatively long term, TCP-friendly flows should not exceed TCP flows in achieved throughput under the same circumstance
- Why TCP-friendly?
 - the majority is still TCP-transported
 - fair competition btw TCP and non-TCP flows
 - important for network stability w/o reservation

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TCP congestion control



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Congestion control mechanisms

- Slow-Start
- Congestion Avoidance
 - increase congestion window size linearly
- Timeout Retransmit
- Fast Retransmit
- Fast Recovery
 - reduce congestion window size to half

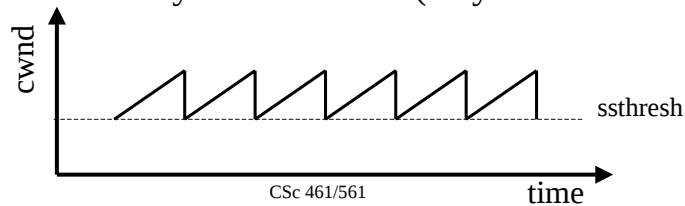
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AIMD

- Additive-Increase-Multiplicative-Decrease
 - why? [hint: fairness and stability]
- TCP: AIMD (1, 0.5)
 - congestion avoidance: increase by 1
 - fast recovery: reduce to 0.5 (why bad for MM?)

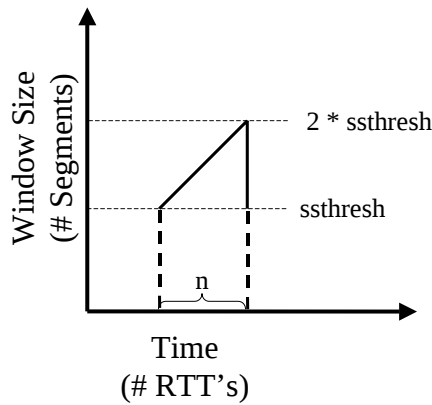


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TCP throughput



- Throughput (T)
 - $\text{data_sent}/\text{time_used}$
- Data transmitted
 - $1.5 n^2 \text{ MSS}$
 - packet loss rate: p
- Time used
 - $n \text{ RTT}$
- $T = \sqrt{1.5/p} \text{ MSS/RTT}$

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Generic AIMD

- GAIMD
 - no congestion: cwnd increased by a ($a > 0$)
 - congestion: cwnd reduced to b ($0 < b < 1$)
- Multimedia GAIMD
 - b close to 1 as much as possible
 - but still responsive to network congestion
 - to reduce instantaneous fluctuation
 - good for smooth multimedia playback

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TCP-friendly GAIMD

- TCP: AIMD (1, 0.5)
- GAIMD (a, b)
- TCP-friendliness
 - $T_{\text{AIMD}} \leq T_{\text{TCP}}$ in a comparatively long term
- Result: $a = 3(1-b)/(1+b)$
 - ref: slide 9 and [CSPM05]
- Examples
 - $(1, 0.5) \sim (0.75, 0.6) \sim (0.53, 0.7) \sim (0.33, 0.8)$

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TCP-friendly CC: more examples

- RAP: rate adaptation protocol
 - apply AIMD to inter-packet gap
- TFRC: TCP-friendly rate control
 - equation-based rate control, with
 - a more sophisticated TCP throughput equation
- Binomial congestion control
 - no congestion: $\text{cwnd}_{i+1} = \text{cwnd}_i + a \text{cwnd}_i^{-x}$
 - congestion: $\text{cwnd}_{i+1} = \text{cwnd}_i - b' \text{cwnd}_i^y$

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This lecture

- TCP-friendly congestion control
 - TCP congestion control
 - AIMD and generic AIMD
 - TCP-friendly congestion control
- Explore further
 - www.psc.edu/networking/projects/tcpfriendly/
 - [CSPM05] Cai, Shen, Pan, Mark, “Performance analysis of TCP-friendly AIMD algorithms for multimedia applications,” *IEEE Trans on Multimedia*, 7(2):339-355, April 2005.

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Next lecture

- Multimedia error control

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