

CSc 360

Operating Systems

Semaphores

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Review: synchronization

- Peterson's solution
 - software-based solution

```
do {  
    flag [i] := true;  
    turn = j;  
    while (flag [j] and turn == j) ;  
    /* critical section */  
    flag [i] = false;  
    /* remainder section */  
} while (1);
```
 - assumption? limitation?

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“Test-and-set”

- Test and set value atomically

```
boolean TestAndSet(boolean &target) {  
    boolean rv = target;  
    target = true;  
    return rv;  
}  
  
boolean lock = false; /* shared variable */  
do {  
    while (TestAndSet(lock)) ;  
    /* critical section */  
    lock = false;  
    /* remainder section */  
}
```

- Any problem?

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Semaphores

- Semaphore API
 - Semaphore S – integer variable
 - binary semaphore
 - counting semaphore
 - two indivisible (atomic) operations
 - also known as P() and V()

```
wait (S):  
    while  $S \leq 0$  do no-op;  
     $S--$ ;
```

```
signal (S):  
     $S++$ ;
```

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Using semaphores

- Mutual exclusion
 - binary semaphore
 - shared data

```
semaphore mutex; //initially mutex = 1
```
 - process P_i

```
do {
    wait(mutex);
    /* critical section */
    signal(mutex);
    /* remainder section */
} while (1);
```
- Resource access
 - counting semaphore
 - initially, the number of resource instances

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Semaphore implementation

- Semaphores without busy waiting
 - block(): block the caller process
 - wakeup(): wakeup another process
- ```
wait(S):
 S.value--;
 if (S.value < 0) {
 add this process to S.L;
 block();
 }

signal(S):
 S.value++;
 if (S.value <= 0) {
 remove a process P from S.L;
 wakeup(P);
 }
```

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## More on using semaphores

- Ordered execution
  - initially, `flag = 0`;
  - P1: ...; *do\_me\_first*; `signal (flag)`;
  - P2: ...; `wait (flag)`; *then\_follow\_on*;
- Caution
  - deadlock
    - `wait (A)`; `wait (B)`; ...; `signal (A)`; `signal (B)`;
    - `wait (B)`; `wait (A)`; ...; `signal (B)`; `signal (A)`;

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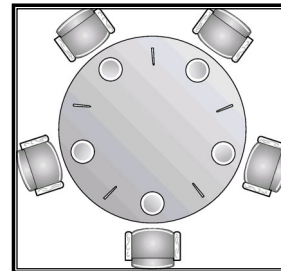
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## Example: dining philosophers

- Shared data
  - Initially all values are 1

```
semaphore chopstick[5];
```
- Philosopher *i*:

```
do {
 wait(chopstick[i])
 wait(chopstick[(i+1) % 5])
 ...
 eat
 ...
 signal(chopstick[i]);
 signal(chopstick[(i+1) % 5]);
 ...
 think
 ...
} while (1);
```



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## This lecture

- Hardware-assisted synchronization
  - test-and-set and *swap*
- Semaphores
  - with(out) busy waiting
- Properties
  - mutual exclusion
  - making process
  - bounded waiting

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## Next lecture

- More on synchronization
  - read OSC7Ch6

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