

## CSC 225 Midterm Exam

June 22, 2001

Recall that you need at least 40% (40/100) in order to write the final exam in this course.

1. [20] Assume  $n = 2^k - 1$  for some integer  $k$ . Solve the following recurrence using repeated substitution:  $T(n) = n + 2T(\frac{n-1}{2})$ ,  $T(7) = 5$ .

2. [20] Prove by induction that your solution to question #1 is correct. Or for part marks [10], apply induction to the point where you realize that your solution to #1 is incorrect, and explain what goes wrong.

The recurrence from Question #1:  $T(n) = n + 2T(\frac{n-1}{2})$ ,  $T(7) = 5$ . You may assume that  $n = 2^k - 1$  for some integer  $k$ .

3. Circle **True** or **False** and justify your answer. **No marks will be given unless there is a correct justification.**

- (a) [5] Let  $a_0, a_1, a_2$ , and  $a_3$  be integers where  $a_i > 0$  for all  $i = 0, 1, 2, 3$ . Then  $p(n) = a_0 + a_1n + a_2n^2 + a_3n^3$  is in  $\Theta(n^3)$ .

**True**

**False**

- (b) [5] An algorithm for sorting  $n$  numbers which is  $O(n \log n)$  in the worst case is always faster than an algorithm which is  $O(n^2)$  in the worst case.

**True**

**False**

- (c) [5] It is possible to sort an array of  $n$  numbers in  $O(n \log n)$  time in the worst case.

True

False

- (d) [5] Let  $f, g$ , and  $h$  be functions from the natural numbers to the positive real numbers. Then if  $g \in \Omega(f)$  and  $g \in O(h)$ , and  $f \in O(h)$  then  $g \in \Theta(h)$ .

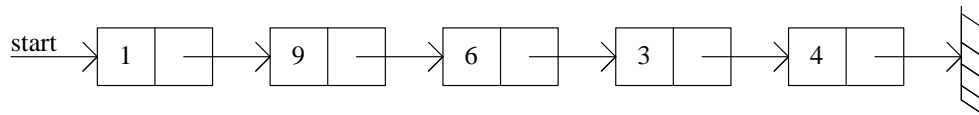
**True**

**False**

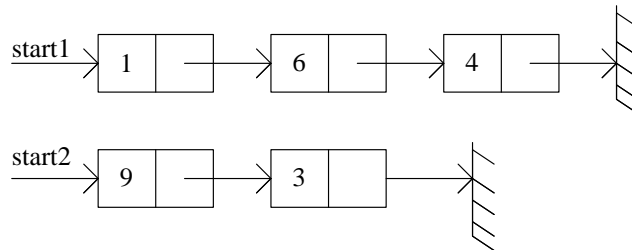
4. [20] Suppose we are given as input two linked lists  $L_1$  (with  $n_1$  nodes, and start/rear pointers  $start_1$  and  $rear_1$ ) and  $L_2$  (with  $n_2$  nodes, and start/rear pointers  $start_2$  and  $rear_2$ ). The objective is to create a new linked list  $L$  with  $n$  nodes, and start/rear pointers  $start$  and  $rear$  such that  $L$  is  $L_1$  concatenated with  $L_2$ . Give detailed pseudocode for an algorithm for this that takes  $O(1)$

time. Marks will be deducted for correct solutions that are more complex than necessary.

5. [20] Give pseudocode for an **iterative** *divide/split* function for merge sort which takes as input a linked list  $L$  starting at  $start$ , and returns as output two lists  $L_1$  and  $L_2$  starting at  $start1$  and  $start2$  respectively. The algorithm should work by placing the first cell from  $L$  on  $L_1$ , then it should place the second cell from  $L$  on  $L_2$ , then the third cell from  $L$  goes on  $L_1$ , and the fourth cell from  $L$  goes on  $L_2$ , and so on. For example, if the input is:



The output should be:



Be sure to include lots of comments in your pseudocode.