

# Prolog II

- Prolog
- Facts (only head)
  - mammal(human) <-
- Query
  - <- mammal(x), legs(x,y)
  - Horn clause without a head

# Resolution and Unification (how queries are expressed)

- $a \leftarrow a_1 \dots a_n$
- $b \leftarrow b_1 \dots b_m$
- If  $b_i$  matches  $a$  then we can infer the clause:
- $b \leftarrow b_1, \dots, b_{i-1}, a_1, \dots, a_n, b_{i+1}, \dots, b_m.$
- Another view: combine left hand / right hand cancel
- $b \leftarrow a.$  and  $c \leftarrow b.$        $b, c \leftarrow a, b$  gives  $c \leftarrow a$

# Resolution

- Goal or list of goals is a Horn clause without a head
- Match one of the goals with the head of known clause
- Simplest case
  - `mammal(human). <- (fact)`
  - `<-mammal(human). (query)`
  - `mammal(human) <- mammal(human)`
  - `<- (query is proved)`

# Unification

- Making two terms “the same”
- $me = me$ 
  - yes
- $me = you$ 
  - no
- $me = X.$ 
  - $X = me$
- $f(a,X) = f(Y,b).$ 
  - $X = b \ Y = a$

# Computation

- Goal: is a a list of goal as a Horn clause without head
- Attempt to apply resolution by matching one of the goal with head of known clause
- Then replace with body, new list of goals
- Repeat until elimination of all goals (proved)

# An example

Facts and rules:

$\text{legs}(x,2) \leftarrow \text{mammal}(x), \text{arms}(x,2).$

$\text{legs}(x,4) \leftarrow \text{mammal}(x), \text{arms}(x,0).$

$\text{mammal}(\text{horse}) \leftarrow.$

$\text{arms}(\text{horse},0) \leftarrow.$

Query:

$\leftarrow \text{legs}(\text{horse},4).$

Resolution:

$\text{legs}(x,4) \leftarrow \text{mammal}(x), \text{arms}(x,0), \text{legs}(\text{horse},4).$

Unification:

$\text{legs}(\text{horse},4) \leftarrow \text{mammal}(\text{horse}), \text{arms}(\text{horse},0), \text{legs}(\text{horse},4)$   
 $\leftarrow \text{mammal}(\text{horse}), \text{arms}(\text{horse},0).$

Resolution

$\text{mammal}(\text{horse}) \leftarrow \text{mammal}(\text{horse}), \text{arms}(\text{horse},0).$

$\leftarrow \text{arms}(\text{horse},0).$

$\text{arms}(\text{horse},0) \leftarrow \text{arms}(\text{horse},0).$

$\leftarrow$

Initial query is true

# Arithmetic

- `write(3+).` evaluates to `3+5`
- `X is 3+5, write(X)` `X = 8`

Gcd in Prolog:

```
gcd(U,0,U).
```

```
gcd(U,V,W) :- R is U mod V, gcd(V,R,W).
```

# Lists

- $[a, b, c]$
- $[a, b, c]$  can also be written  $[a, b, c | []]$  or  $[a, b | [c]]$  or  $[a | [b, c]]$
- $[H|T] = [a, b, c]$ 
  - $H = a, T = [b, c]$
- $[a|T] = [H, b, c]$ 
  - $T = [b, c], H = a$
- $[H, T]$  is syntactic sugar for  $.(H, T)$  ( $.$  is cons)



# Actual code examples

- ancestor
- links
- append

# Important

Queries are yes/ fail rather than yes/ no

No means the system can not prove it, not that is necessarily false

Prolog:

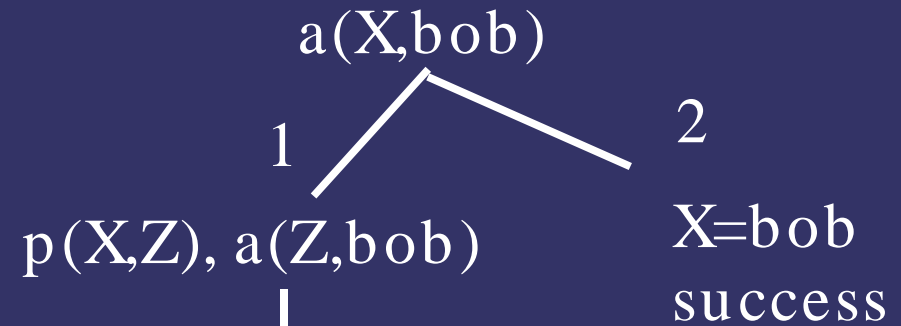
Order of clauses top-to-bottom

Order of goals left-to-right

Always depth-first search

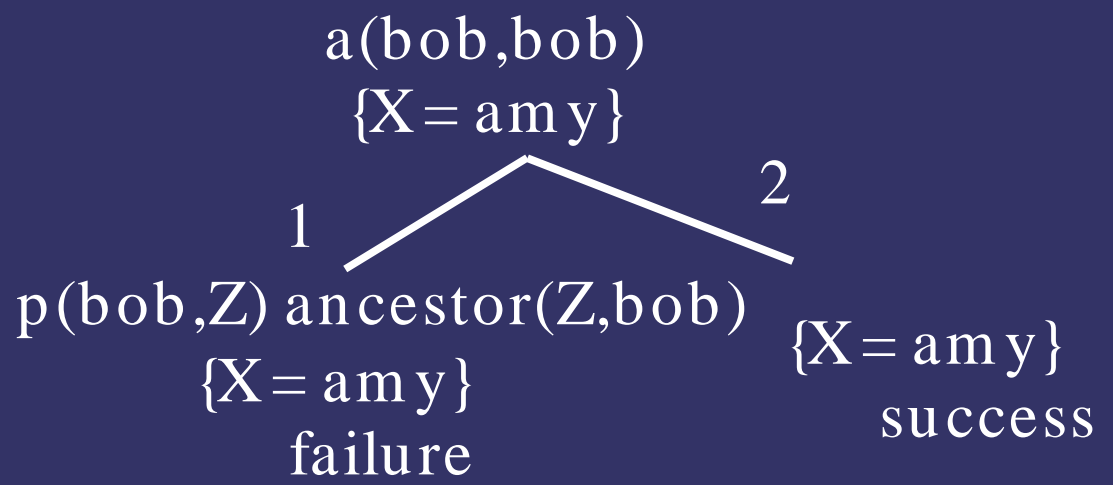
# Prolog Search Tree

- 1 a(X,Y) :- p(X,Z), a(Z,Y).
- 2 a(X,X).
- 3 p(amy,bob).



Depth-first search strategy

Problem:  
 a(X,Y) :- a(Z,Y), p(X,Z)  
 goes into an infinite loop



ORDER MATTERS

# Cuts

- 1  $a(X,Y) :- p(X,Z), !, a(Z,Y).$
- 2  $a(X,X).$
- 3  $p(\text{amy},\text{bob}).$

Cut “freezes” the choice made, if it is reached on backtracking, the subtrees of parent node are not examined. Cut “prunes” the search tree.

Can be used for efficiency

