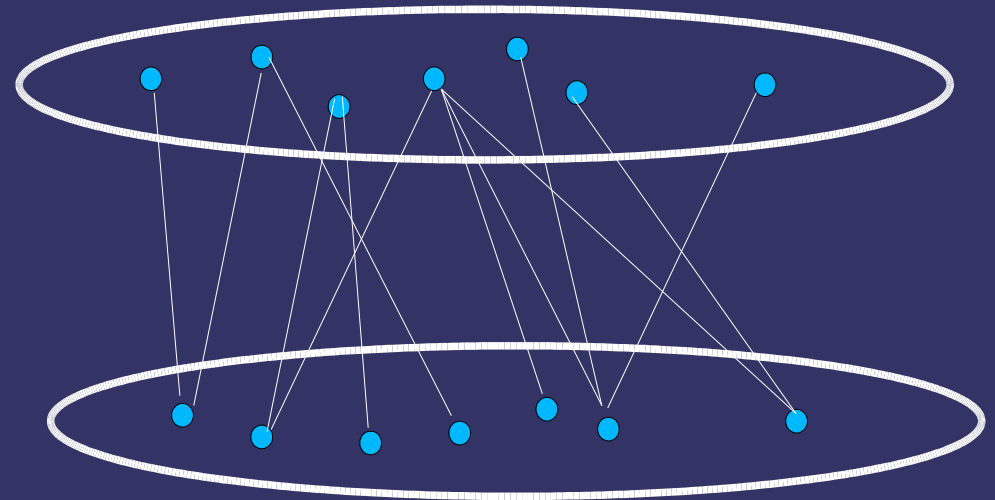
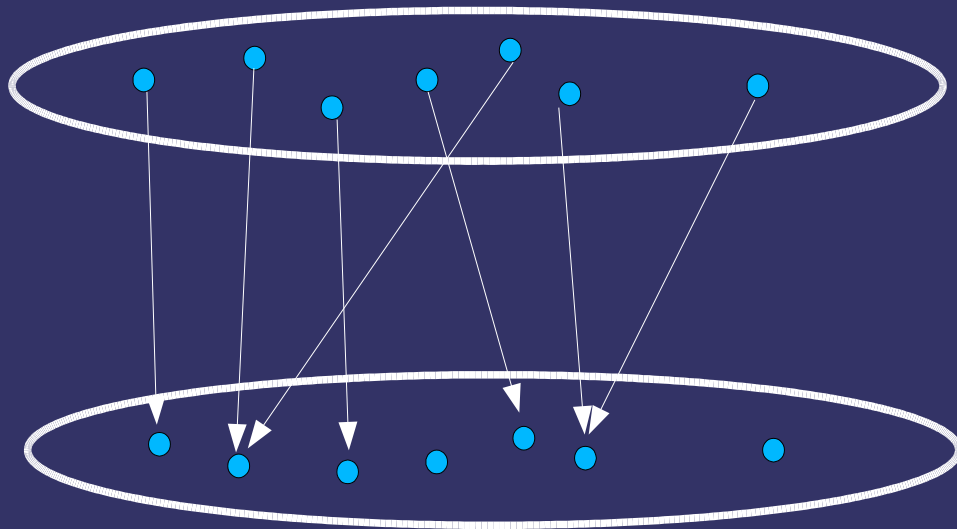


Prolog

- Logic programming
 - Origins: automatic deduction systems, theorem provers
 - Basic idea: computation can be viewed as a kind of proof
- Prolog (1970s)
 - 1981 Japan's fifth generation project

Overview

- Programs in functional and imperative languages are mappings (many to one)
- Logic programs are relations (many to many)



Append

- Relation append is a set of tuples of form (X, Y, Z) where Z consists of the elements of X followed by the elements of Y .

$([a],[b],[a,b])$ is in relation append

$([a],[b], [])$ is not in relation append

First-order predicate calculus

- Constants : numbers/names
- Predicates : functions that are true or false
- Functions : non-boolean values
- Variables : unspecified quantities
- Connectives : and, or, not, implication ->
- Quantifiers : for all, there exists

Logical statements

In English:

A horse is a mammal

A human is a mammal

Mammals have four legs and no arms, or two legs and two arms

A horse has no arms

A human has arms

In FOPC:

mammal(horse).

mammal(human).

for all x, mammal(x) \rightarrow

legs(x,4) and arms(x,0) or legs(x,2) and arms(x,2)

arms(horse,0).

not arms(human,0).

Inference rule

- Infer: legs(horse,4).
- Axioms, theorems proved by inference

$(a \rightarrow b) \text{ and } (b \rightarrow c)$

 $a \rightarrow c$

A logical programming language is a notational system for writing logical statements together with specific algorithms for implementing inference rules

How does it work ?

Facts:

mammal(horse).

mammal(human).

for all x, mammal(x) ->

legs(x,4) and arms(x,0) or legs(x,2) and arms(x,2)

arms(horse,0).

not arms(human,0).

Query: there exists y, legs(human, y) ?

Answer: yes: $y = 2$

Deductive:

Specify properties of solution
and find it without specifying
exactly how

Horn Clauses

- Horn clauses
 - a_1 and a_2 and a_3 and ... $a_n \rightarrow b$
 - body implies head
- Can express most, but not all, logical statements

An example

English: x is a grandparent of y if x is the parent of someone who is the parent of y .

First-order predicate calculus:

for all x , for all y , (there exists z , $\text{parent}(x,z)$ and $\text{parent}(z,y)$)
 $\rightarrow \text{grandparent}(x,y)$.

Horn clause:

$\text{parent}(x,z)$ and $\text{parent}(z,y) \rightarrow \text{grandparent}(x,y)$

Procedural interpretation

- $b \leftarrow a_1 \text{ and } a_2 \text{ and } a_3 \dots \text{ and } a_n$
 - viewed as a procedure for obtaining b
- $\text{sort}(x,y) \leftarrow \text{permutations}(x,y) \text{ and } \text{sorted}(y)$

$\text{gcd}(u,0,u)$.

$\text{gcd}(u,v,w) \leftarrow \text{not zero}(v), \text{gcd}(v, u \bmod v, w)$.

Resolution and Unification (how queries are expressed)

- $a \leftarrow a_1 \dots a_n$
- $b \leftarrow b_1 \dots b_m$
- If b_i matches a then we can infer the clause:
- $b \leftarrow b_1, \dots, b_{i-1}, a_1, \dots, a_n, b_{i+1}, \dots, b_m.$

An example

Facts and rules:

legs(x,2) <- mammal(x), arms(x,2).

legs(x,4) <- mammal(x), arms(x,0).

mammal(horse).

arms(horse,0).

Query:

<- legs(horse,4).

Resolution:

legs(x,4) <- mammal(x), arms(x,0), legs(horse,4).

Unification:

legs(horse,4) <- mammal(horse), arms(horse,0), legs(horse,4)
<- mammal(horse), arms(horse,0).

Resolution

mammal(horse) <- mammal(horse), arms(horse,0).

<- arms(horse,0).

arms(horse,0) <- arms(horse,0).

<-

Initial query is true

Prolog

ISO Prolog based on Edinburgh Prolog (de facto standard today)

```
ancestor(X,Y) :- parent(X,Z), ancestor(Z,Y).  
ancestor(X,X).  
parent(amy,bob).
```

Order can be important:
ancestor(x,bob).

If left to right then x is amy
If right to left then x is bob

Actual code example

Queries

Queries are yes/ fail rather than yes/ no
No means I can not prove it