

CS 330 Lecture 14

> Outline

- > Call-by-value and Call-by-need (lazy evaluation)
- > Tail recursion
- > Efficiency
- > Infinite data structures

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Two simple functions

- > `fun sqr(x) = x * x`
- > `fun zero(x) = 0;`
- > When a function is called, the argument is substituted for the function's formal parameter in the body
- > Evaluation rule: expressions have multiple function calls – when and how many times is the argument evaluated

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Evaluation in ML: call-by-value

- > Simple expressions: constants, variables, function calls and conditional expressions
- > ML's evaluation rule:
 - > To compute the value of $f(E)$, first compute the value of expression E then substitute in f
- > `sqr(sqr(sqr(2))) -> sqr(sqr(2x2)) -> sqr(sqr(4)) -> sqr(4x4) -> sqr(16) -> 16 * 16 => 256`
- > what about `zero(sqr(sqr(sqr(2))))` ?

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Evaluation of recursive functions

- > `fun fact = if n = 0 then 1 else n * fact(n-1)`

```
fact(4) -> 4 * fact(4-1)
         -> 4 * fact(3)
         -> 4 * (3 * fact(3-1))
         -> 4 * (3 * (2 * fact(2-1)))
         -> 4 * (3 * (2 * (1 * fact(1-1))))
         -> 4 * (3 * (2 * (1 * fact(0))))
         -> 4 * (3 * (2 * (1 * 1)))
         -> 4 * (3 * (2 * 1))
         -> 4 * (3 * 2)
         -> 4 * 6
         -> 24
```

This seems inefficient.
The larger number the more numbers wait to be multiplied.
Can we do better ?

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Tail-recursive factorial

- › Associative law $4 * (3 * (\text{fact}(2))) = 12 * \text{fact}(2)$
- › fun facti (n,p) = if n=0 then p else facti(n-1, n * p)
- › tail recursive = the result of recursive call is returned immediately without modification

```
facti(4,1) -> facti(4-1, 4x1)
           -> facti(3, 4)
           -> facti(2, 12)
           -> facti(2-1, 2 * 12)
           -> facti(1-1, 1 * 24)
           -> facti(0, 24)
           -> 24
```

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Special role of conditional expressions

- › Conditional expressions (if – then – else, andalso, orelse) are not functions
- › fun cond(p,x,y) = if p then x else y;
- › fun badf n = cond(n=0, 1, n*badf(n-1));
- › Every call to badf runs forever why ?

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Call-by-name

- › What if we give functions their arguments as expressions, not values ?
- › Call-by-name rule: to compute the value of $f(E)$, substitute E immediately into the body of f . Then compute the value of the resulting expression
- › $\text{zero}(\text{sqr}(\text{sqr}(\text{sqr}(2))))$ is immediately 0
- › However observe what happens in:
 $\text{sqr}(\text{sqr}(\text{sqr}(2))) = \text{sqr}(\text{sqr}(2)) * \text{sqr}(\text{sqr}(2))$

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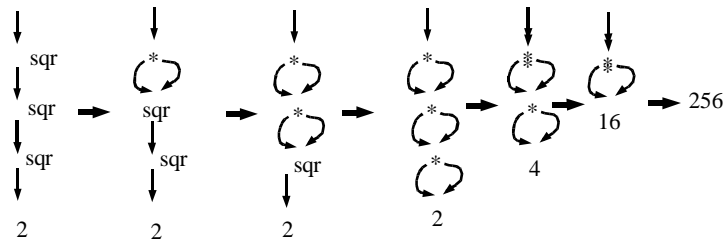
Call-by-need – Lazy evaluation

- › Like call-by-name but ensure that each argument is evaluated at most once
- › Pointer structure – Directed graph of functions and arguments
- › Graph reduction

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Graph reduction of $\text{sqr}(\text{sqr}(\text{sqr } 2))$



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Comparison of strict and lazy evaluation

- › Call-by-need does the least possible evaluation but requires much book-keeping
- › Sometimes lazy evaluation saves a lot of spaces; sometimes it wastes space

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Space leak with lazy evaluation for facti

```
facti(4,i) -> facti(4-1, 4 * 1)
            -> facti(3-1, 3 * (4 * 1))
            -> facti(2-1, 2 * (3 * (4 * 1)))
            -> facti(1-1, 1 * (2 * (3 * (4 * 1))))
            -> 1 * (2 * (3 * (4 * 1)))
...
            -> 24
```

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Haskell

- › www.haskell.org
- › Pure functional language with lazy evaluation
- › Lazy lists = elements are not evaluated until their values are needed by the rest of the program
- › All data structures in Haskell are Lazy
- › As an exercise let's do infinite lists in SML

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