

Lecture 29

- > Prolog
- > Facts (only head)
 - > mammal(human) <-
- > Query
 - > <- mammal(x), legs(x,y)
 - > Horn clause without a head

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Resolution and Unification (how queries are expressed)

- > $a \leftarrow a_1 \dots a_n$
- > $b \leftarrow b_1 \dots b_m$
- > If b_i matches a then we can infer the clause:
 - > $b \leftarrow b_1, \dots, b_{i-1}, a, \dots, a_n, b_{i+1}, \dots, b_m$.
- > Another view: combine left hand /right hand cancel
- > $b \leftarrow a.$ and $c \leftarrow b.$ $b, c \leftarrow a, b$ gives $c \leftarrow a$

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Resolution

- > Goal or list of goals is a Horn clause without a head
- > Match one of the goals with the head of known clause
- > Simplest case
 - > mammal(human). <- (fact)
 - > <-mammal(human). (query)
 - > mammal(human) <- mammal(human)
 - > <- (query is proved)

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Unification

- > Making two terms "the same"
- > $me = me$
 - > yes
- > $me = you$
 - > no
- > $me = X.$
 - > $X = me$
- > $f(a,X) = f(Y,b).$
 - > $X = b$ $Y = a$

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Computation

- > Goal: is a list of goal as a Horn clause without head
- > Attempt to apply resolution by matching one of the goal with head of known clause
- > Then replace with body, new list of goals
- > Repeat until elimination of all goals (proved)

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An example

Facts and rules:

legs(x,2) <- mammal(x), arms(x,2).

legs(x,4) <- mammal(x), arms(x,0).

mammal(horse) <-.

arms(horse,0) <-.

Query:

<- legs(horse,4).

Resolution:

legs(x,4) <- mammal(x), arms(x,0), legs(horse,4).

Unification:

legs(horse,4) <- mammal(horse), arms(horse,0), legs(horse,4)

<- mammal(horse), arms(horse,0).

Resolution

mammal(horse) <- mammal(horse), arms(horse,0).

<- arms(horse,0).

arms(horse,0) <- arms(horse,0).

<-

Initial query is true

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Prolog

ISO Prolog based on Edinburgh Prolog (de facto standard today)

ancestor(X,Y) :- parent(X,Z), ancestor(Z,Y).

ancestor(X,X).

parent(amy,bob).

Order can be important:

ancestor(x,bob).

If left to right then x is amy

If right to left then x is bob

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Arithmetic

> write(3+). evaluates to 3+5

> X is 3+5, write(X) X = 8

Gcd in Prolog:

gcd(U,O,U).

gcd(U,V,W) :- R is U mod V, gcd(V,R,W).

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Lists

- > [a, b, c]
- > [a,b,c] can also be written [a,b,c | []] or [a, b | [c]] or [a | [b, c]]
- > [H|T] = [a,b,c]
 - > H = a, T = [b,c]
- > [a|T] = [H,b,c]
 - > T = [b,c], H = a
- > [H,T] is syntactic sugar for `.(H,T)` (`.` is cons)

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Actual code examples

- > ancestor
- > links
- > append

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Important

Queries are yes/fail rather than yes/no
No means the system can not prove it, not that is necessarily false

Prolog:

Order of clauses top-to-bottom
Order of goals left-to-right

Always depth-first search

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Prolog Search Tree

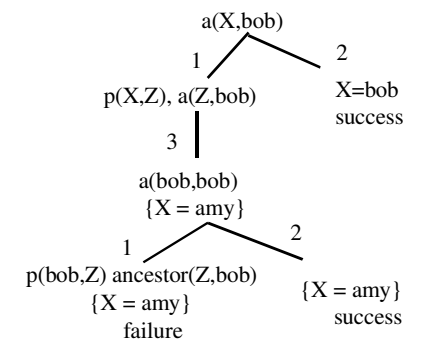
1 a(X,Y) :- p(X,Z), a(Z,Y).
2 a(X,X).
3 p(amy,bob).

Depth-first search strategy

Problem:

a(X,Y) :- a(Z,Y), p(X,Z)
goes into an infinite loop

ORDER MATTERS



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Cuts

- 1 $a(X,Y) :- p(X,Z), !, a(Z,Y).$
- 2 $a(X,X).$
- 3 $p(\text{amy},\text{bob}).$

Cut “freezes” the choice made, if it is reached on backtracking, the subtrees of parent node are not examined. Cut “prunes” the search tree.

Can be used for efficiency

